

What are Reactive rule-based systems?

- Agents whose behavior is defined by Reactive rules
 - E-commerce platforms (react to user actions, e.g. put an item in the basket)
 - Web services (react to notifications, e.g. SOAP messages)
 - Active databases

Reactive rules

- Reactive rules families:
 - Production rules (PR)
 - Event Condition Action rules (ECA)
 - Knowledge representation rules (KR)
 - Complex Event Processing (CEP)



Reactive rules

- Bridge the **gap** between existing and dynamic Web
- Used to specify reactive systems
- Support ad-hoc, flexible,
 dynamic workflows that change at run-time

- No formal verification tools
- Interaction of rules during execution
- Unpredictable behavior of rules



Problem

- Not enough appropriate formal tools
 - for **intelligent agents** with complex behavior
- Addressing their characteristics consistently is not trivial
- Need for developing methodologies targeted to those systems
- Meet their challenges and requirements

Formal methods

- Unambiguous description of a system (Specification)
- Effective analysis of desired properties (Verification)
- Algebraic specification languages
 - Combinations of logical systems
- CafeOBJ
 - Specification of complex systems as abstract state machines
 - Verification of safety properties through theorem proving



Why Formal methods?

- Testing and simulation **not enough** to guarantee system's behavior
- Better understanding of the specified system
- Reasoning support
- A system does not have defects
- A system does **satisfy** desirable **properties**
- Constructing **reliable** distributed systems



In a nutshell...

- Need for verifying structure and safety properties about reactive rules
 - -Exploit CafeOBJ's support for **rewriting logic** specifications
- Express reactive rules as rewrite rules

In a nutshell...

- PR and Rewriting logic / PR and Equational Logic in CafeOBJ
- ECA and Rewriting logic / ECA and Equational Logic in CafeOBJ
- Proving termination properties using rewriting logic approach
- Proving confluence properties using rewriting logic approach
- Simulating rules' execution behavior using rewriting logic approach
- Proving safety properties using both approaches

A light-control intelligent system

- ECA rules specifying a light-control intelligent system
- Energy consumption reduction
 - Turning off the lights in unoccupied rooms or when the occupant is asleep using sensors
- Automatic adjustment for indoor light intensity based on the outdoor light intensity



environme- ntal variables	Mtn, ExtLgt, Slp				
local variables	lgtsTmr, intLgts				
external	MtnOn activated when $Mtn = true$				
events	MtnOff activated when $Mtn = false$				
	ExtLgtLow activated when $ExtLgt \leq 5$				
internal events	SecElp, LgtsOff, LgtsOn, ChkExtLgt, ChkMtn, ChkSlp				
(R1)	When the room is unoccupied for 6 minutes, turn off lights if they are on.				
rl	on MtnOff if (intLgts > 0 and lgtsTmr $= 0$) do set (lgtsTmr, 1) par activate (SecElp)				
r2	on SecElp if (lgtsTmr ≥ 1 and lgtsTmr < 6 and Mtn = false) do increase (lgtsTmr, 1)				
r3	on SecElp if (lgtsTmr = 6 and Mtn = false) do set (lgtsTmr, 0) par activate (LgtsOff)				
r4	on LgtsOff do (set (intLgts, 0) par activate (ChkExtLgt))				
(R2)	When lights are off, if external light intensity is below 6, turn on lights.				
r5	on ChkExtLgt if (intLgts = 0 and ExtLgt ≤ 5) do activate (LgtsOn)				
(R3)	When lights are on, if the room is empty or a person is asleep, turn off lights.				
r6	on LgtsOn do (set (intLgts, 6) seq activate (ChkMtn))				
r7	on ChkMtn if (Slp = true or (Mtn = false and intLgts >= 1)) do activate (LgtsOff)				
(R4)	If the external light intensity drops below 5, set the lights intensity to 6 and check if the person is asleep. If the person is asleep, turn off the lights.				
r8	on ExtLgtLow do set (intLgts, 6) par activate (ChkSlp)				
r9	on (ChkSlp if $(Slp = true)$ do set (intLgts, 0)				
(R5)	If the room is occupied, set the lights intensity to 4.				
r10	on MtnOn do set (intLgts, 4) par set (lgtsTmr, 0)				

Environmental variables: store the values measured by sensors

Mtn: motion sensor detecting whether the room is occupied or not

Slp: pressure sensor detecting whether the person is asleep or not

ExtLgt: light sensor for monitoring the outdoor lighting

MtnOn, MtnOff, ExtLgtLow: external events (activated by environmental variables)

Internal events: internal actions

Desired safety property:

- The lights cannot be turned off if someone is in the room and he/she does not sleep
- inv1(S) = (not((intLgts(S) = 0) and (Mtn(S) = true) and (Slp(S) = false))



Comparison of methodologies

Equational approach

- More expressive formalism
- Preferred when the reactive rule-based system is complex

Rewriting approach

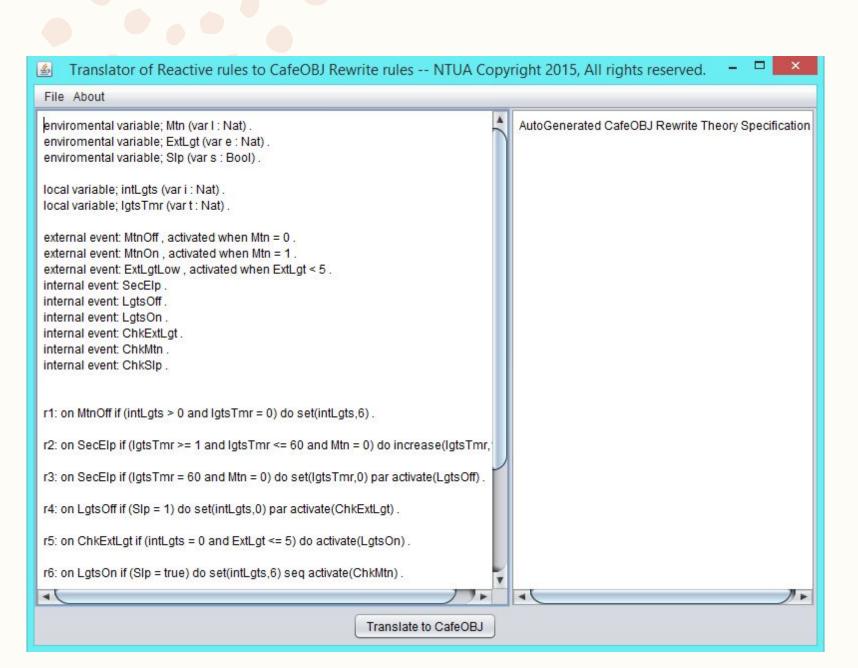
- More natural and easier to use from reactive rules researchers
- Seamless framework for verifying both safety properties and structure errors
- Supports both theorem proving and model checking techniques
- **Preferred** in most cases



Tool: from reactive to rewrite rules

- Transforms a set of reactive rules into a CafeOBJ rewrite theory specification
- Written in Java
- Syntactic guidelines of reactive rules specification:
 - Rules declaration should start with an identification number (e.g. r1)
 - Definition of variables, events and rules should end with a fullstop (.)
 - External events should be accompanied with their detection conditions





ECA rules of the light control intelligent agent



About	
omental variable; Mtn (var I : Nat) .	
omental variable; ExtLgt (var e : Nat).	mod! RULES {
omental variable; SIp (var s : Bool) .	pr(STATE)
	pr(EQL)
variable; intLgts (var i : Nat).	op null : -> Name .
variable; lgtsTmr (var t : Nat).	
	var n : Name
al event: MtnOff , activated when Mtn = 0 .	
nal event: MtnOn , activated when Mtn = 1 .	var I : Nat
al event: ExtLgtLow , activated when ExtLgt	var e : Nat
al event: SecElp .	var s : Bool
al event: LgtsOff.	
al event: LgtsOn .	vari: Nat
al event: ChkExtLgt .	vart:Nat
al event: ChkMtn .	
al event: ChkSlp .	ops MtnOff MtnOn ExtLgtLow SecElp LgtsOff LgtsOn ChkExtLgt ChkMtn Cl
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	ctrans [A1]: (event-memory: MtnOff) (lgtsTmr: t) => (event-memory: null) (intLgts
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SecElp if (lgtsTmr >= 1 and lgtsTmr <= 60	strong (A2); (sugget memory OppElp) ((staTopy t) (()the I) -> (sugget memory Late
Section if (late Tran - 60 and Mtn - 0) do not	<pre>ctrans [A3]: (event-memory: SecElp) (lgtsTmr: t) (Mtn: l) => (event-memory: Lgts</pre>
SecElp if (IgtsTmr = 60 and Mtn = 0) do set	ctrans [A4]: (event-memory: LgtsOff) (intLgts: i) (SIp: s) => (event-memory: ChkE
LgtsOff if (SIp = 1) do set(intLgts,0) par acti	
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ChkExtLgt if (intLgts = 0 and ExtLgt <= 5) do	
ormexiegt in (integro = o and extegr = 5) d	ctrans [A6a]: (event-memory: LgtsOn) (intLgts: i) (SIp: s) => (event-memory: Lgts
LgtsOn if (SIp = true) do set(intLgts,6) seq	

Click on **Translate to CafeOBJ** button



Franslator of Reactive rul	es to CafeOBJ Rewrite rules	NTUA Copyright 2015, All rights reserved.	-
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Additional CafeOBJ specification The output of the tool can be given directly as input to the CafeOBJ processor

Future work

- Integrate reactive rules verification methodology with other rewriting systems
 - e.g. Maude \rightarrow better tool support and broader use



Research Interests

- Formal methods
 - Algebraic specifications
 - Automated theorem provers
- Intelligent Software Agents
- Complex Systems
 - Blockchains



